

## Homework 3

(150 points)

1. **(15 + 15 + 15 + 15 points) Fourier Analysis on Larger Domains.** Recall that we apply discrete Fourier Analysis on the Boolean Hypercube to analyze functions with domain  $\{0, 1\}^n$ . We will generalize this construction to arbitrary domains.

- (a) Consider the space of all function  $\mathbb{Z}_p \rightarrow \mathbb{C}$ , where  $p$  is a prime number. Here  $\mathbb{Z}_p$  is the set  $\{0, 1, \dots, p-1\}$ . And addition and multiplication of two elements from this set is defined using integer addition and multiplication, respectively,  $\pmod p$ . The set of complex numbers is represented by  $\mathbb{C}$ .

Suppose  $f, g: \mathbb{Z}_p \rightarrow \mathbb{C}$  be two functions. Recall that the *complex conjugate* of a complex number  $z = a + ib$ , represented by  $\bar{z}$ , is defined to be  $a - ib$ . The inner-product of these two functions is defined by

$$\langle f, g \rangle := \frac{1}{p} \sum_{x \in \mathbb{Z}_p} f(x) \overline{g(x)}$$

Let  $\omega_p := \exp(2\pi i/p)$  and define  $\chi_a(x) := \omega_p^{ax}$ , for  $a \in \mathbb{Z}_p$ . Prove that  $\{\chi_a : a \in \mathbb{Z}_p\}$  is an orthonormal basis for the space of all function  $\mathbb{Z}_p \rightarrow \mathbb{C}$ .

- (b) Consider the space of all functions  $\mathbb{Z}_p^n \rightarrow \mathbb{C}$ . Define the inner-product of functions, write the Fourier basis functions, and show their orthonormality.
- (c) Consider the space of all functions  $\mathbb{Z}_p \times \mathbb{Z}_q \rightarrow \mathbb{C}$ , for primes  $p$  and  $q$ . The primes  $p$  and  $q$  need not necessarily be distinct. Define the inner-product of functions, write the Fourier basis functions, and show their orthonormality.
- (d) Consider the space of all functions  $\mathbb{Z}_{p_1} \times \mathbb{Z}_{p_2} \times \dots \times \mathbb{Z}_{p_n} \rightarrow \mathbb{C}$ . Note that the primes  $p_1, \dots, p_n$  need not be distinct. Define the inner-product of functions, write the Fourier basis functions, and show their orthonormality.



2. **(10 + 10 points) Majority Functions.** Let  $n$  be odd and  $f(x): \{0, 1\}^n \rightarrow \{+1, -1\}$  be the majority function. That is, if the majority of the bits in  $x$  is 0, then  $f(x) = +1$ ; otherwise  $f(x) = -1$ .

(a) Compute the Fourier coefficients of  $f$  when  $n = 3$ .

(b) Let us define odd and even functions. For  $x \in \{0, 1\}^n$ , define  $\text{flip}(x)$  to be the string where we flip every bit of  $x$ . For example  $\text{flip}(0010) = 1101$ .

A function is *odd* if  $f(\text{flip}(x)) = -f(x)$ , for all  $x \in \{0, 1\}^n$ . Note that the majority function defined above is an odd function.

A set  $S \in \{0, 1\}^n$  is *even* if the number of 1s in  $S$  is even. For example, when  $n = 3$ , the sets  $S = 000, 011, 101, 110$  are even sets.

Prove that if  $f$  is an odd function then  $\widehat{f}(S) = 0$  for all even  $S \in \{0, 1\}^n$ .



3. **(20 points) Generalized BLR.** Recall that a function  $f: \{0,1\}^n \rightarrow \{+1, -1\}$  is linear if  $f(0^n) = +1$  and  $f(x + y) = f(x) \cdot f(y)$ , for all  $x, y \in \{0,1\}^n$ . Consider the following generalization of the BLR algorithm to test whether a function  $f$  is close to linear or the function  $-f$  is close to linear.

BLR – Gen<sup>f</sup>:

- (a) Let  $a, b, c \xleftarrow{\$} \{0,1\}^n$
- (b) Let  $w = f(a)$ ,  $x = f(b)$ ,  $y = f(c)$ , and  $z = f(a + b + c)$
- (c) Return  $(w \cdot x \cdot y == z)$

State and prove a theorem that intuitively proves that “the algorithm returns true with high probability” if and only if “the function  $f$  or  $-f$  is close to a linear function.”



4. **(20 points) An Alternate Proof.** Recall that the convolution of two function  $f, g: \{0, 1\}^n \rightarrow \mathbb{R}$  is defined as follows

$$(f * g)(x) := \frac{1}{N} \sum_{y \in \{0,1\}^n} f(y)g(x - y)$$

In this problem we shall develop a new technique to prove that  $\widehat{(f * g)} = \widehat{f}\widehat{g}$ .

- (a) Compute the function  $(\chi_S * \chi_T)$
- (b) Note that the convolution operator is a bilinear operator. That is, we have  $((f_1 + f_2) * g) = (f_1 * g) + (f_2 * g)$  and  $(cf) * g = c(f * g)$  from the definition of convolution. Similarly, we have  $(f * (g_1 + g_2)) = (f * g_1) + (f * g_2)$  and  $f * (cg) = c(f * g)$ .

Recall that we have  $f = \sum_{S \in \{0,1\}^n} \widehat{f}(S)\chi_S$  and  $g = \sum_{S \in \{0,1\}^n} \widehat{g}(S)\chi_S$ . Prove that

$$(f * g) = \sum_{S \in \{0,1\}^n} \widehat{f}(S)\widehat{g}(S)\chi_S$$





5. (5 + 15 + 5 + 5 points) **A Few Properties of Fourier Transformation.** Let  $f, g: \{0, 1\}^n \rightarrow \mathbb{R}$  be two functions.

- (a) Express  $\widehat{fg}$  using the functions  $\widehat{f}$  and  $\widehat{g}$ . Here the function  $(fg)$  defined as  $(fg)(x) = f(x) \cdot g(x)$ , for all  $x \in \{0, 1\}^n$ .
- (b) Let  $\max\{f, g\}$  is the function that satisfies  $\max\{f, g\}(x) = \max\{f(x), g(x)\}$ , for all  $x \in \{0, 1\}^n$ . Suppose the range of  $f$  and  $g$  is  $\{+1, -1\}$ . Express  $\widehat{\max\{f, g\}}$  in terms of  $\widehat{f}$  and  $\widehat{g}$ .
- (c) Recall that if  $f(x) = g(x - c)$  for some  $c \in \{0, 1\}^n$  then we have  $\widehat{f} = \chi_c \widehat{g}$ . Find a function  $h: \{0, 1\}^n \rightarrow \mathbb{R}$  such that  $f = (h * g)$ .
- (d) For  $1 \leq i < j \leq n$ , define

$$\text{swap}_{i,j}(x_1, \dots, x_n) = (x_1, \dots, x_{i-1}, x_j, x_{i+1}, \dots, x_{j-1}, x_i, x_{j+1}, \dots, x_n)$$

. Suppose  $f(x) = g(\text{swap}_{i,j}(x))$ , for all  $x \in \{0, 1\}^n$ . Express  $\widehat{f}$  as a function of  $\widehat{g}$ .